

Minimax lower bounds (continued)

Chapter 11

Recall

$$\begin{array}{l} H_0 \quad p_0 \sim \gamma \\ \underline{vs} \quad H_1 \quad p_1 \sim \gamma \end{array}$$

$$\rho = \sum_y \sqrt{p_0(y) p_1(y)}$$

Lemma

$$\frac{p^2}{4} \leq p_{e^*} \leq \frac{p}{2}$$

↳ minimum average probability of error

Example

Example

$$Y = (Y_1, \dots, Y_n) \sim \begin{matrix} \text{Be}(\frac{1-h}{2}) & H_0 \\ \text{Be}(\frac{1+h}{2}) & H_1 \end{matrix} \quad |h| \geq h \geq 0$$

$\longleftrightarrow h$

$$\left(\begin{array}{l} n=1 \\ p_0 = \left(\frac{1+h}{2}, \frac{1-h}{2} \right) \\ p_1 = \left(\frac{1-h}{2}, \frac{1+h}{2} \right) \end{array} \right. \quad \rho = 2 \sqrt{\frac{1-h^2}{4}} = \sqrt{1-h^2}$$

A number line with arrows at both ends. Three points are marked with dots and labeled below: $\frac{1-h}{2}$, $\frac{1}{2}$, and $\frac{1+h}{2}$.

$$n \geq 1 \quad \rho = (\sqrt{1-h^2})^n \geq 0$$

$$\rho^2 = (1-h^2)^n$$

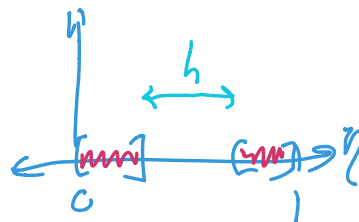
$$p^2 = (1-h^2)^n \quad \frac{(1-h^2)^n}{4} \leq p_e^* \leq \frac{(1-h^2)^{n/2}}{2}$$

$$\left(h = \frac{1}{\sqrt{n}} \right) \quad (1-h^2)^n = \left(1 - \frac{1}{n} \right)^n \approx e^{-1}$$

$P(h)$ - $\mathcal{P}$ such that for any x ,

if $\eta(x) = P\{Y=1 | X=x\}$, then

$$\eta(x) \in [0, \frac{1-h}{2}] \cup [\frac{1+h}{2}, 1]$$



Thm 12.1 Let $\mathcal{F}$ have VC dimension

$V \geq 2$, Then for any $n \geq V$, and

$$h \in [0, 1]$$

$$R_n(h, \mathcal{F}) \geq c \min\left(\sqrt{\frac{V}{n}}, \frac{V}{nh}\right) \quad \left(c = \frac{1}{32} \text{ works}\right)$$

$\uparrow$ min max expected excess risk

12.2 $X, \mathcal{F}: X \rightarrow \{0, 1\}, P \in \mathcal{P}(h)$

$$L(f) = P\{Y \neq f(X)\}$$

$$f^*(x) = 1_{\{\eta(x) \geq 1/2\}} \quad (\text{MAP estimator})$$

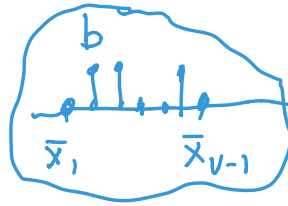
$$\underline{R_n(h, \mathcal{F})} \triangleq \inf_A \sup_{P \in \mathcal{P}(h)} E[L(\hat{f}_n) - L^*] \quad \begin{matrix} (\min \\ \max) \end{matrix}$$

$$\hat{f}_n = A(Z^n)$$

$$Z^n = ((X_1, Y_1), \dots, (X_n, Y_n))$$

Proof Let $\bar{x}_1, \dots, \bar{x}_V$ be V points in X that are shattered by $\mathcal{F}$.

For any $b \in \{0, 1\}^{V-1}$ there exists f_b & g such that $f_b(\bar{x}_v) = b_v$ $1 \leq v \leq V-1$
 (let $b_v = 0$) $f_b(\bar{x}_v) = 0$



Use flip probability $\frac{1-h}{2}$.

Let P_b denote the joint distⁿ on $X \times \{0, 1\}$ so

if (X, Y) has distⁿ P_b then

$$X = \begin{cases} \bar{x}_v & \text{w/p } p \frac{1}{V-1} \text{ if } 1 \leq v \leq V-1 \\ \bar{x}_V & \text{" } 1-p \end{cases} \quad b_v$$

$$\text{Given } X = \bar{x}_v, \quad P_b(Y=1 | X=\bar{x}_v) = \frac{1-h}{2} + h \mathbb{1}_{\{b_v=1\}}$$

Define 2^{V-1} probability distributions on $X \times \{0, 1\}$.

Let $\bar{P}^n$ be the joint distⁿ of (B, Z^n) where $Z^n = ((X_1, Y_1), \dots, (X_n, Y_n))$ as follows.

Let B be uniformly distributed over $\{0, 1\}^{V-1}$

Given B , let $Z^n | B=b$ have distⁿ P_b .

Let $\hat{f}_n$ be output of an arbitrary algorithm.

$$L_B(\hat{f}_n) = \frac{1-h}{2} + \frac{ph}{V-1} \sum_{v=1}^{V-1} \mathbb{1}_{\{\hat{f}_n(\bar{x}_v) \neq B_v\}}$$

$$E^n[L_B(\hat{f}_n)] - \frac{1-h}{2} = \frac{ph}{V-1} \sum_{v=1}^{V-1} P\{\hat{f}_n(\bar{x}_v) \neq B_v\}$$

expected excess risk if B_v is true distⁿ per sample

$$\frac{1}{2^{V-1}} \sum_{b \in \{0,1\}^{V-1}} \left[E^n[L_b(\hat{f}_n)] - \frac{1-h}{2} \right] = \frac{ph}{V-1} \sum_{v=1}^{V-1} P\{\hat{f}_n(\bar{x}_v) \neq b_v\}$$

$$\therefore R_n(h, \mathbb{F}) \geq \sup_b \left[E^n[L_b(\hat{f}_n)] - \frac{1-h}{2} \right] \geq ph \left(\frac{1}{V-1} \sum_{v=1}^{V-1} P\{\hat{f}_n(\bar{x}_v) \neq b_v\} \right)$$

Lower bound.

Let $N_v = \#$ of X_i 's equal to $\bar{x}_v$. $1 \leq v \leq V-1$

By Bhattacharyya bound

$$\bar{P}^n\{\hat{f}_n(\bar{x}_v) \neq B_v | N_v\} \geq \frac{1}{4}(1-h^2)^{N_v}$$

$$E[N_v] = \frac{np}{V-1} \quad N_v \stackrel{d}{=} u_1 + \dots + u_n \quad u_i \sim \text{Ber}\left(\frac{p}{V-1}\right)$$

$$\text{So } \bar{P}^n\{\hat{f}_n(\bar{x}_v) \neq B_v\} = E\left[\right] =$$

$$= \frac{1}{4} \left(\frac{p}{V-1} (1-h^2) + \frac{1-p}{V-1} \right)^n$$

$$= \frac{1}{4} \left(1 - \frac{ph^2}{V-1} \right)^n$$

$$\boxed{\text{So } (*) R_n(h, f) \geq \frac{ph}{4} \left(1 - \frac{ph^2}{v-1}\right)^n} \quad \text{all } \begin{array}{l} n \geq 1 \\ v \geq 2 \\ h \in [a, 1] \\ p \in [a, 1] \end{array}$$

Case 1 $\sqrt{\frac{v-1}{n}} \leq h \leq 1$

Let $p = \frac{v-1}{(n+1)h^2}$. Then

$$(*) \text{ yields } R_n(h, f) \geq \frac{v-1}{4(n+1)h} \left(1 - \frac{1}{n+1}\right)^n \geq \frac{v-1}{16hn} \\ \geq \frac{v}{32hn}$$

Case 2 $0 \leq h \leq \sqrt{\frac{v-1}{n}}$ — use bound for case 1 with $h = \sqrt{\frac{v-1}{n}}$.